## Characterizing Truthfulness In Discrete Domains

AHUVA MU'ALEM

Social and Information Sciences Laboratory, California Institute of Technology and

MICHAEL SCHAPIRA

The School of Computer Science and Engineering, The Hebrew University of Jerusalem,  $\mathsf{Israel}^1$ 

Algorithmic mechanism design [9; 10] focuses on the design of algorithms that aim to achieve global objectives in settings in which the "input" is provided by self-interested strategic players<sup>2</sup>. This necessitates the design of algorithms that are *incentive-compatible* (a.k.a. *truthful*<sup>3</sup>) in the sense that players are incentivized via payments to behave as instructed. The most natural approach to designing incentive-compatible algorithms is coming up with an algorithm *and* an explicit payment scheme that guarantees its incentive-compatibility. However, finding appropriate payments is often a difficult, setting-specific, task, which is mostly achievable for very simple types of algorithms.

A more general approach is the following: Any algorithm that interacts with selfish players and then outputs an outcome, can be regarded as computing a function, called a *social-choice function*, from the players' "input" to some outcome space. Certain properties of social-choice functions are known to imply their *imple-mentability*, that is, the *existence* of a payment scheme that guarantees incentive-compatibility. Hence, instead of explicitly dealing with payments, the problem of designing incentive-compatible algorithms boils down to analyzing the mathematical properties of the social-choice functions computed by algorithms. This approach makes sense if these mathematical properties are simple and easy to analyze.

A simple constraint on social choice functions called "*weak-monotonicity*" has been shown to characterize the implementability of social choice functions in several interesting settings. However, with the exception of very restricted settings named "single parameter domains" [1], all these characterizations of incentive-

<sup>&</sup>lt;sup>1</sup>Supported by grants from the Israel Science Foundation.

 $<sup>^2\</sup>mathrm{We}$  deal with the standard quasilinear mechanism design setting.

<sup>&</sup>lt;sup>3</sup>Formally, the solution concept we consider is incentive-compatibility in ex-post Nash (in particular, our results hold for incentive-compatibility in dominant strategies, which is a special case of this solution concept).

Authors' addresses: ahumu@caltech.edu, mikesch@cs.huji.ac.il

Permission to make digital/hard copy of all or part of this material without fee for personal or classroom use provided that the copies are not made or distributed for profit or commercial advantage, the ACM copyright/server notice, the title of the publication, and its date appear, and notice is given that copying is by permission of the ACM, Inc. To copy otherwise, to republish, to post on servers, or to redistribute to lists requires prior specific permission and/or a fee. © 2008 ACM 1529-3785/2008/0700-0001 \$5.00

## 2 • A. Mu'alem and M. Schapira

compatibility are known to apply only to environments in which the private information of the players is drawn from inherently non-discrete domains, like convex domains (see [2; 4; 12; 8]). Our work [13] is motivated by the fact that in many cases the private information of the players is drawn from discrete domains (e.g., integers). Implementability in discrete domains is still little understood and has received but little attention in economic literature [7].

We consider the following standard mechanism design setting: There are n players  $1 \ldots n$ , and a set of outcomes O. Each player i has a private valuation function  $v_i \in V_i$  that assigns a real value to every  $o \in O$  (the higher the value of the outcome the more desirable it is). A (deterministic) social-choice function is a function that assigns an outcome o to every  $v \in V$ , where V denotes  $V_1 \times \ldots \times V_n$ . Let  $V_{-j}$  denote the cartesian product of all  $V_i$ s but  $V_j$ , and let  $(v_i, v_{-i})$  denote the profile of valuation functions in which player i's valuation function is  $v_i \in V_i$ , and the other players' valuation functions are as specified by  $v_{-i} \in V_{-i}$ . Then, f is implementable iff there is a payment function  $p_i$  such that for every  $i \in [n]$ , for every  $v_{-i} \in V_{-i}$ , and for every  $v_i, v'_i \in V_i$ ,

$$v_i(f(v_i, v_{-i})) - p_i(v_i, v_{-i}) \ge v_i(f(v_i', v_{-i})) - p_i(v_i', v_{-i}).$$

Rochet [11] has shown that any social-choice function is implementable iff a constraint called "cycle monotonicity" holds (see recent application by Lavi and Swamy [6]). Bikhchandani et al. [2] proposed the simple weak-monotonicity constraint: A social-choice function f is said to be weakly-monotone if for every  $i \in [n]$ , for every  $v_{-i} \in V_{-i}$ , and for every  $v_i, v'_i \in V_i$ , such that  $f(v_i, v_{-i}) = o_1$  and  $f(v'_i, v_{-i}) = o_2$  it holds that

$$v_i(o_1) + v'_i(o_2) \ge v_i(o_2) + v'_i(o_1).$$

f is said to be strongly-monotone if whenever  $o_1 \neq o_2$  this inequality is strict [5]. It is easy to show that weak monotonicity is always necessary for the implementability of a social-choice function but that it is not always sufficient.

We start our exploration of incentive-compatibility by focusing on two important types of discrete domains: Integer grid domains and 0/1 domains [13]. Let  $V = V_1 \times \ldots \times V_n$  be a domain of valuation functions defined over a set of outcomes O. We can think of every  $v_i \in V_i$  as a vector in  $\mathbb{R}^{|O|}$  specifying a value for every outcome.

DEFINITION 1. A valuation function domain is an integer grid domain if  $V = Z^{|O|} \times \ldots \times Z^{|O|}$ .

That is, an integer grid domain is a domain of valuation functions that can take any combination of integer values. We exhibit an example (due to Lan Yu) that shows that in integer grids weak monotonicity is *not* sufficient to guarantee implementability [13]. In fact, this example can easily be made to hold for any *bounded* integer grid. By bounded integer grid, we simply mean the discrete cube  $V = \{0, 1, ..., L\}^{n|O|}$  (for some positive integer  $L \ge 1$ ). In contrast, we show that strong monotonicity is sufficient to obtain implementability in integer grids.

DEFINITION 2.  $V = V_1 \times ... \times V_n$  is a 0/1 domain if  $V = \{0, 1\}^{|O|} \times ... \times \{0, 1\}^{|O|}$ . ACM SIGecom Exchanges, Vol. 7, No. 2, June 2008.

3

We show that, as in the case of integer grids, in 0/1 domains weak-monotonicity is insufficient for implementability, but strong-monotonicity is.

OPEN QUESTION 1. Is strong-monotonicity sufficient for implementability in bounded integer grids?

When does weak monotonicity guarantee implementability in discrete domains? As we have seen this is not true even in natural discrete settings (like integer grids). In contrast, we present [13] a family of discrete domains in which weak-monotonicity suffices for implementability, which we term *Monge domains*. The proof that this is indeed true for Monge domains takes advantage of the two dimensional version of submodularity (see [3]) that holds for this kind of domains (expressed by Monge matrices, hence the name). Monge domains have a simple combinatorial structure that has many advantages from a mechanism design perspective. We refer the reader to [13] for more details.

## REFERENCES

Aaron Archer and Eva Tardos. Truthful mechanisms for one-parameter agents. In FOCS, pages 482–491, 2001.

S. Bikhchandani, S. Chatterji, R. Lavi, A. Mu'alem, N. Nisan, and A. Sen. Weak monotonicity characterizes deterministic dominant strategy implementation. *Econometrica*, 74(4):1109–1132, July 2006.

V.G. Deineko and G. Woeginger. Some problems around travelling salesmen, dart boards, and euro-coins. *Bulletin of the European Association for Theoretical Computer Science*, 90:43–52, October 2006.

Hongwei Gui, Rudolf Muller, and Rakesh Vohra. Characterizing dominant strategy mechanisms with multi-dimensional types, 2004. Working paper.

Ron Lavi, Ahuva Mu'alem, and Noam Nisan. Towards a characterization of truthful combinatorial auctions. In FOCS, 2003.

Ron Lavi and Chaitanya Swamy. Truthful mechanism design for multi-dimensional scheduling via cycle monotonicity. In EC, 2007.

William S. Lovejoy. Optimal mechanisms with finite agent types. *Manage. Sci.*, 52(5):788–803, 2006.

Dov Monderer. Monotonicity and implementability, 2007. Working paper.

Noam Nisan and Amir Ronen. Algorithmic mechanism design. *Games and Economic Behavior*, 35:166–196, 2001.

Noam Nisan, Tim Roughgarden, Eva Tardos, and Vijay V. Vazirani (eds.). Algorithmic Game Theory. Cambridge University Press, 2007.

J.C. Rochet. A necessary and sufficient condition for rationalizability in a quasi-linear context. *Journal of Mathematical Economics*, 16:191–200, 1987.

Michael Saks and Lan Yu. Weak monotonicity suffices for truthfulness on convex domains. In EC, 2005.

Ahuva Mu'alem and Michael Schapira. Mechanism design over discrete domains. In EC, 2008.