Auctions with Correlated and Interdependent Values

Part I: Contributions of Robert Wilson and Paul Milgrom

Presented by Inbal Talgam-Cohen (Technion)
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OF ALFRED NOBEL 2020

Paul R. Milgrom Robert B. Wilson

"for improvements to auction theory and inventions of new auction formats"

Foundational Papers We'll Touch Upon

Analysis of leading auction formats (1st price, 2nd price, English)

- 1. Wilson'77, "A bidding model of perfect competition", The Review of Economic Studies
- Milgrom'81, "Rational expectations, information acquisition, and competitive bidding", Econometrica
- Milgrom-Weber'82, "A theory of auctions and competitive bidding", Econometrica

Auction design and robustness

- Cremer-McLean'88, "Full extraction of the surplus in Bayesian and dominant strategy auctions", Econometrica
- 5. Wilson'87, "Game-theoretic analyses of trading processes", Advances in Economic Theory: Fifth World Congress

Let's Revisit the Common-Value Setting

Setting:

- Symmetric bidders
- Every bidder *i*:
 - has value $u_i(S, X) = S$
 - knows only her own signal X_i
- Signal distribution is known

Notation:

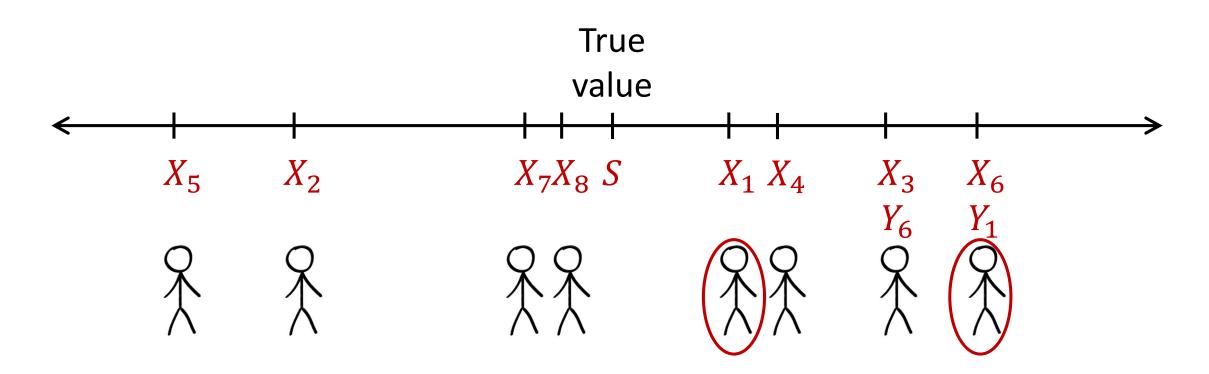
• Y_i = highest signal of bidders other than i



Value S

Common-Value Setting

• For concreteness let $X_i = S + \epsilon_i$ where ϵ_i 's are i.i.d. and $\mathbb{E}[\epsilon_i] = 0$



Three Values of Interest

Initial estimate of S by bidder i:

$$\hat{S}_i = \mathbb{E}[S \mid X_i]$$

- Estimate of S upon winning the auction:
 - $\widehat{W}_i = \mathbb{E}[S \mid X_i, i \text{ won, everything learned during auction}]$
- Bidder *i*'s bid $\beta(x)$
 - where bidding strategy $\beta(\cdot)$ is symmetric, increasing

\hat{S}_i = value estimate \hat{W}_i = estimate upon win $\beta(x)$ = bid

Is Winning Good or Bad News?

Let's first consider 1^{st} price auctions, with bids according to strategy β

• A priori, i's estimate \hat{S}_i of S is simply $X_i = S + \epsilon_i$, an unbiased estimator:

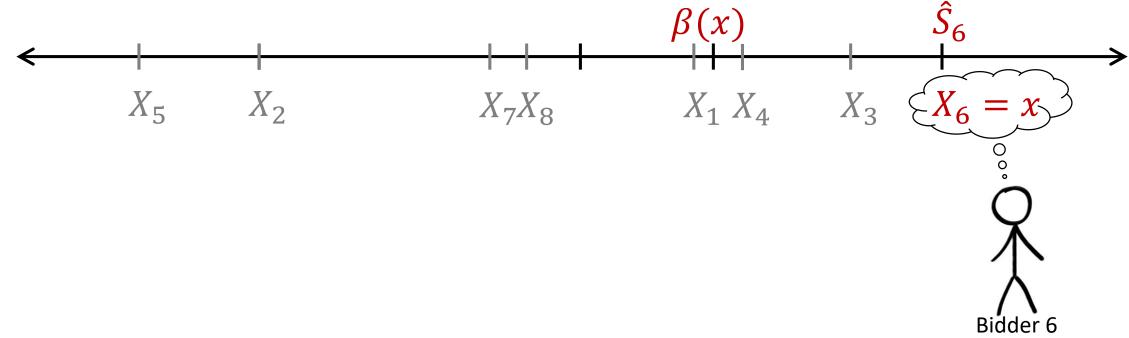
$$\mathbb{E}[X_i \mid S = s] = s$$

• Upon winning, i learns $X_i = \max_{i'}\{X_{i'}\}$, but now: $\mathbb{E} \Big[\max_{i'}\{X_{i'}\} \mid S = s \Big] > s$

Conclusion:
$$\widehat{W}_i < \widehat{S}_i$$
, i.e., winning is bad news

Winner's Curse

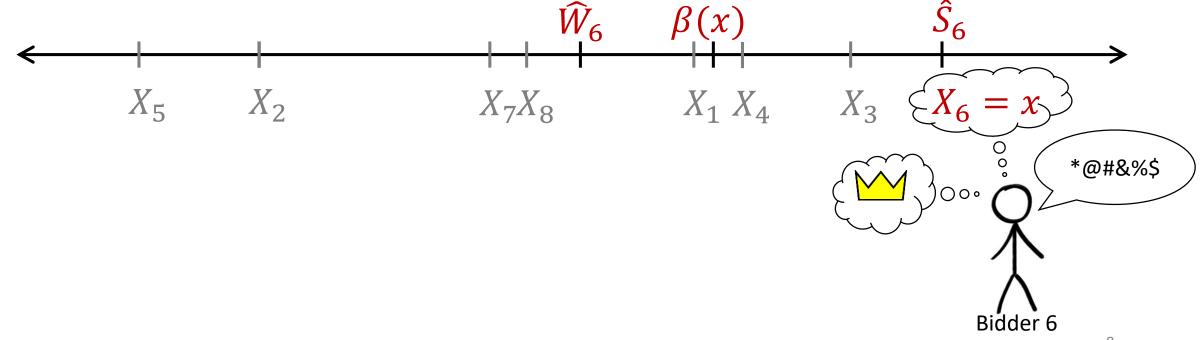
 \hat{S}_i = value estimate \hat{W}_i = estimate upon win $\beta(x)$ = bid



Winner's Curse

$$\hat{S}_i$$
 = value estimate
 \hat{W}_i = estimate upon win
 $\beta(x)$ = bid

• Winner's curse occurs when $\beta(x)$ does not anticipate the shift from \hat{S}_i to \hat{W}_i



Winner's Curse: Implication to 1st Price

- Let $n \to \infty$
 - Enough information on the market to estimate S accurately

 But if the winner's curse occurs, the price converges to something higher than S

Conclusion: Winner's curse \Rightarrow price fails to aggregate information

Does winner's curse occur in equilibrium?

Contribution of [Wilson'77]

- Derives closed-form equilibrium bidding strategy $\beta(\cdot)$
 - $\beta(\cdot)$ is symmetric, increasing under appropriate regularity conditions (affiliation)
- In equilibrium, bidders shade their bids to avoid the winner's curse:

$$\beta(x) \le \widehat{W}_i = \mathbb{E}[S \mid X_i = x, Y_i \le x]$$

$$X_i \text{ is highest}$$

Conclusion: No winner's curse in equilibrium; price converges to true value!

What about 2nd Price Auctions?

Insight from 1st price analysis:

Equilibrium bids avoid the winner's curse by factoring in the anticipated learning from winning

• Let's use this insight to find equilibrium $\beta(\cdot)$ for 2^{nd} price

Learning from Winning in 2nd Price

• Consider 2^{nd} price auctions, with bids according to strategy β

- What does i learn if she wins?
 - $X_i \geq Y_i$
 - $\beta(Y_i)$
- So $\widehat{W}_i = \mathbb{E}[S \mid X_i = x, Y_i = y]$
- (Cf., $\widehat{W}_i = \mathbb{E}[S \mid X_i = x, Y_i \leq x]$ in 1st price)

Equilibrium Strategy [Milgrom'81]

• No knowledge of $\widehat{W}_i = \mathbb{E}[S \mid X_i = x, Y_i = y]$ at bid time...

<u>Theorem</u>: In 2nd price auctions, under regularity conditions, the following bidding strategy is a symmetric increasing equilibrium:

$$\beta(x) = \mathbb{E}[S \mid X_i = x, Y_i = x]$$

• Indeed, winner *i* has no incentive to change bid after learning price:

Affiliated Signals

If some signals are high, the remaining signals are more likely to be high

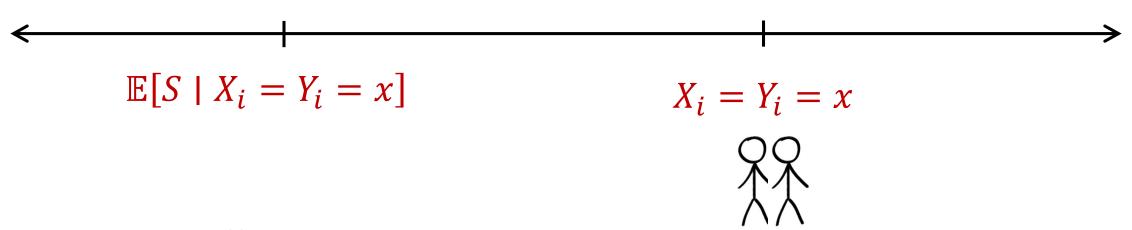
<u>Def</u>: Random variables $X_1, ..., X_n$ with joint density f are affiliated if for every two realizations $x, y \in \mathbb{R}^n$, $f(x)f(y) \leq f(x \wedge y)f(x \vee y)$ Component Component -wise min -wise max

- Equivalent to log-supermodularity of f
- E.g., $f(2,1,8)f(9,7,1) \le f(2,1,1)f(9,7,8)$

Affiliation Implies Increasing Bid $\beta(\cdot)$

$$\beta(x) = \mathbb{E}[S \mid X_i = Y_i = x]$$

 S, X_i, Y_i affiliated



• Intuition: Affiliation means positive correlation everywhere

Linkage Principle [Milgrom-Weber'82]

Overview of Contributions

Consider a general symmetric model:

$$u_i(S, \mathbf{X}) = u(X_i, \{X_j\}_{j \neq i})$$
 for all i

Regularity assumption: affiliation

Results:

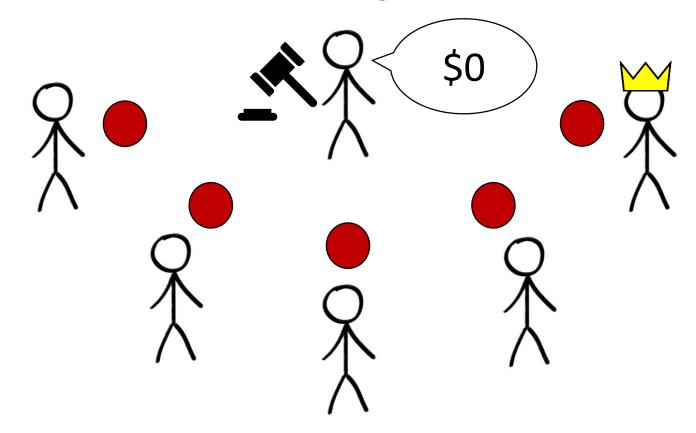
- 1. Symmetric increasing equilibria for English, 1st, 2nd price auctions
- 2. Revenue hierarchy:

$$Rev[1^{st} price] \le Rev[2^{nd} price] \le Rev[English]$$

3. Hierarchy explained by a general Linkage Principle

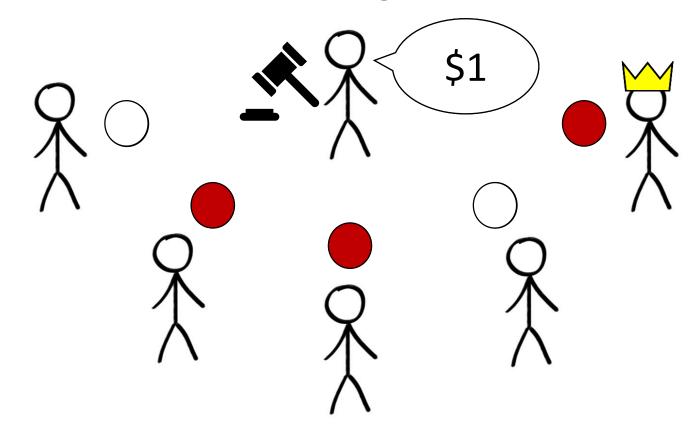
Recall: Equil. strategy factors in the anticipated learning from winning

What does i learn if she wins the English auction?



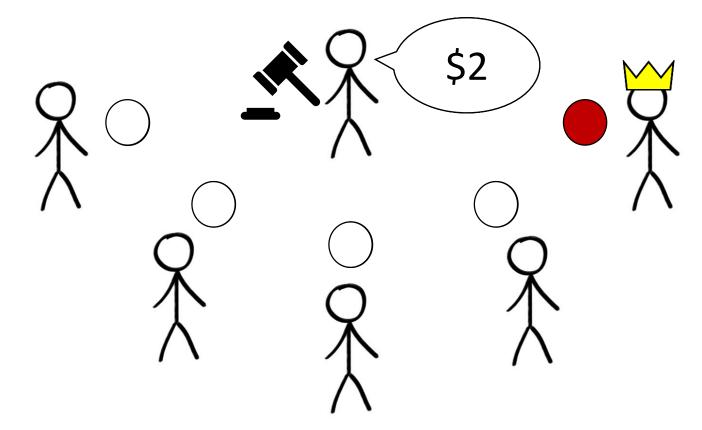
Recall: Equil. strategy factors in the anticipated learning from winning

What does i learn if she wins the English auction?



Recall: Equil. strategy factors in the anticipated learning from winning

What does i learn if she wins the English auction?



Recall: Equil. strategy factors in the anticipated learning from winning

- What does i learn if she wins the English auction?
- Claim: All signals $\{X_i\}$!
 - Since all bidders use the same strategy to decide to drop out, when a bidder drops out her signal is revealed

Equilibrium: Bidding $u(X_i, \{X_j\}_{j\neq i})$ is impossible, so start with $u(x, \{x, ..., x\})$ and at each stage plug in the revealed X_j 's

The Linkage Principal

- A = auction in (symmetric, increasing) equilibrium
- $P^{A}(b, x) =$ equilibrium price if *i* wins with bid *b* and signal *x*
- $P^{A}(x)$ = statistical linkage of equilibrium price to i's signal x

$$\frac{\partial}{\partial x'} P^A(b, x') \mid_{(b, x') = (x, x)}$$

• For example, in 1st price auctions, $P^{first}(x) = 0$

Theorem [Informal]: For auctions A, B, if $P^A(x) \ge P^B(x) \ \forall x$ then $\text{Rev}[A] \ge \text{Rev}[B]$

The Linkage Principal: Intuition

Linkage of equil. price to winner's signal x

• Theorem [Informal]: For auctions A, B, if $P_2^A(x) \ge P_2^B(x) \ \forall x$ then $Rev[A] \ge Rev[B]$

<u>Intuition</u>: Stronger linkage of the (public) price to the winner's private signal x lowers the information rent

• A.k.a. "publicity effect" [Milgrom'04]

Practical Take-Aways

Popularity of English auctions

- "Honesty is the best policy"
 - If seller has information affiliated with that of bidders', revealing it increases revenue

Motivation for auction design

Auction Design & Robustness

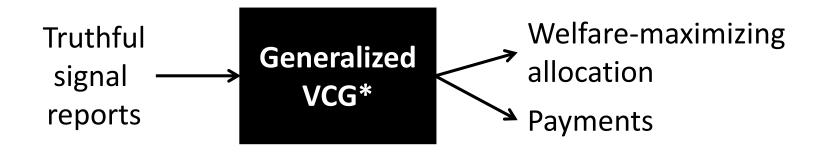
Optimal Auction [Cremer-McLean'88]

- We've seen: Correlation brings down information rents
- [Cremer-McLean'88] takes this to the extreme down to zero!

Main result: A truthful optimal mechanism that extracts the full welfare as revenue

• Generically, even the slightest degree of correlation among signals suffices

Optimal Mechanism: High-Level Idea



- Let $p_i(x_i)$ be the expected payoff (utility) of i from Gen-VCG given x_i
- Impose an additional "lottery" payment $\ell(x_{-i})$ on i such that: $\mathbb{E}_{x_{-i}|x_i}[\ell(x_{-i})] = p_i(x_i)$
- Generically, such a lottery exists iff signals are correlated

Downside and Wilson's Doctrine

The optimal mechanism crucially depends on common knowledge of the joint signal distribution

"I foresee the progress of game theory as depending on successive reductions in the base of common knowledge required to conduct useful analyses of practical problems. Only by repeated weakening of common knowledge assumptions will the theory approximate reality"

[Wilson'87]

Summary So Far

Correlated and interdependent values ≠ **IPV**

- Bidders learn about their value from the auction (format matters!)
- Equilibrium bids anticipate this learning
- Correlation lowers information rents (possibly to zero!)

More after the break...

Recommended Further Reading

- Paul Milgrom, Putting Auction Theory to Work, Cambridge University Press, 2004
 - Chapter 5

- Vijay Krishna, Auction Theory, Elsevier, 2010
 - Chapters 6-10