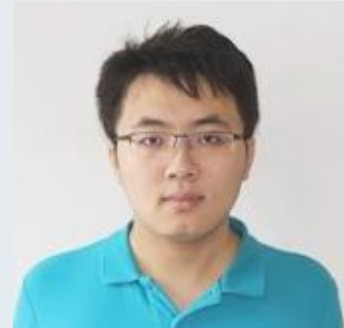


MENU-SIZE COMPLEXITY AND REVENUE CONTINUITY OF BUY-MANY MECHANISMS



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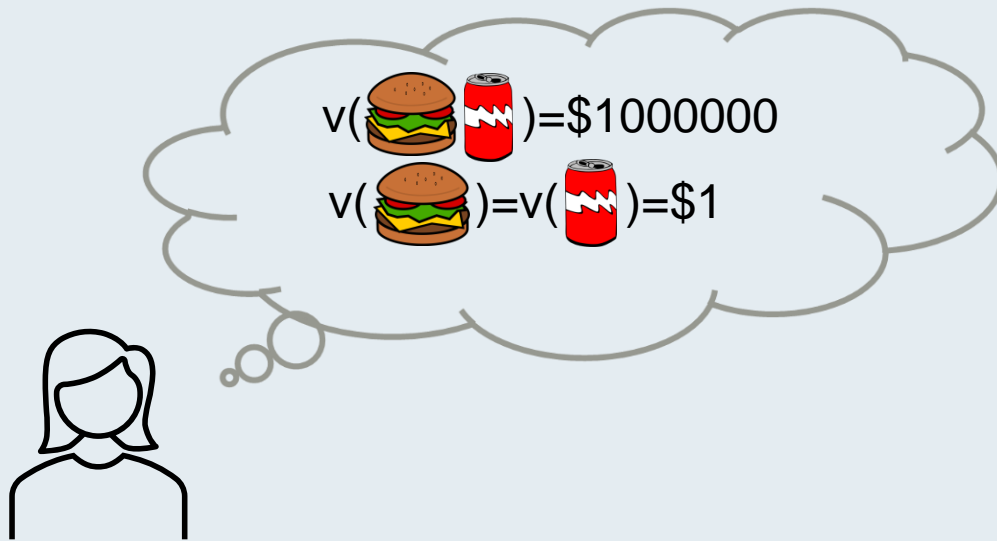
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Buy-one mechanisms and buy-many mechanisms



MENU		
$\frac{1}{2}$		\$5
$\frac{1}{3}$		\$2
		\$999

- A seller has n heterogeneous items to sell to a single buyer.
- Typical buy-one mechanisms: buyer interact with the seller once.
 - *Optimal strategy: purchases the third menu option, pay \$999.*
- Buy-many mechanisms: buyer interact with the mechanism multiple times.
 - *Optimal strategy: repeatedly purchase , then repeatedly purchase , pay \$16 in expectation.*

Menu-size complexity for near-optimal revenue

- How many menu options are needed for $(1 - \epsilon)$ -approx in revenue?
 - *Buy-one mechanisms: infinite [Hart Nisan'13].*
 - *Buy-many mechanisms: finite.*

Theorem 1. For any distribution D and $\epsilon \in [0,1]$, exists mechanism M with finite menu size $f(n, \epsilon)$, such that

$$\text{Rev}_D(M) \geq (1 - \epsilon)\text{BuyManyRev}_D.$$

- $f(n, \epsilon) = (1/\epsilon)^{2^{O(n)}}$.
- The doubly-exponential dependency of n is tight.

Theorem 2. There exists D being a distribution over XOS functions, such that for any mechanism M with description complexity $2^{2^{o(n^{1/4})}}$,

$$\text{BuyManyRev}_D \geq o(\log n)\text{Rev}_D(M).$$

Revenue Continuity

- When the buyer's values for the sets of items perturb multiplicatively slightly, how much does the revenue change?
 - Any $v \sim D$ is perturbed to $v' \sim D'$, such that $v'(S) \in [(1 - \epsilon)v(S), (1 + \epsilon)v(S)]$, $\forall S \subseteq [n]$.
- Buy-one mechanisms: revenue may change significantly [Psomas et al.'19].
 - Continuity only holds for weaker additive perturbation [Rubinstein Weinberg'15] [Brustle et al.'20].
- Buy-many mechanisms: revenue changes slightly.

Theorem 3. For any value distribution D and any $1 \pm \epsilon$ multiplicative perturbation D' ,

$$\text{BuyManyRev}_{D'} \geq (1 - \text{poly}(n, \epsilon)) \text{BuyManyRev}_D.$$

- Note: such dependency on n is necessary.

Theorem 4. There exists D over unit-demand functions and a $1 \pm \epsilon$ multiplicative perturbation D' , such that

$$\text{BuyManyRev}_{D'} \leq \frac{1}{\epsilon n} \text{BuyManyRev}_D.$$

- Full paper: <https://arxiv.org/abs/2003.10636>